



Propel

Verification of Algebraic Properties

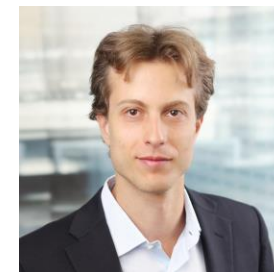
George Zakhour



Pascal Weisenburger



Guido Salvaneschi





Algebraic Properties Are Important!

- Associative
- Commutative
- Idempotent
- Identity
- Zero



Strange Loop
Oct 14-15, 2010
<https://thestrangeloop.com>



Algebraic Properties are Everywhere

- Datastructure Invariants
- Compiler Optimization
- Stream Processing
- Algorithm Design
- Typeclass Laws
- Databases

Algebraic Properties are Everywhere

- Datastructure Invariants
 - a. CRDTs
 - state-based (PLDI' 23)
 - op-based (PLF' 23)
- Compiler Optimization
- Stream Processing
- Algorithm Design
- Typeclass Laws
- Databases



Propel

Type-Checking CRDT Convergence

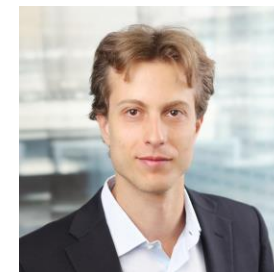
George Zakhour



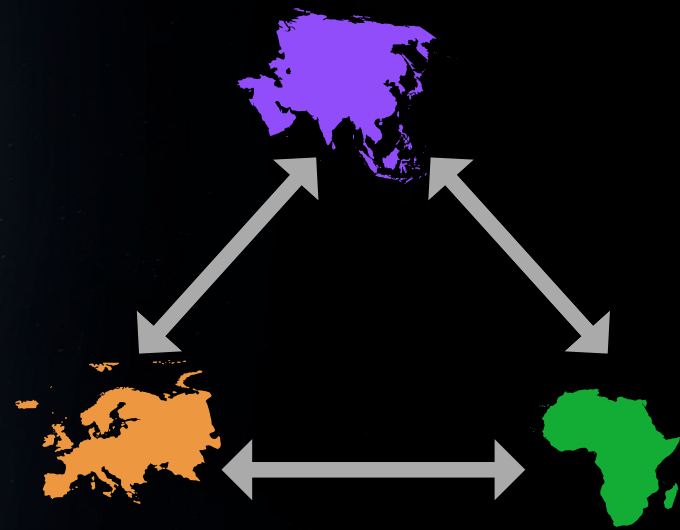
Pascal Weisenburger



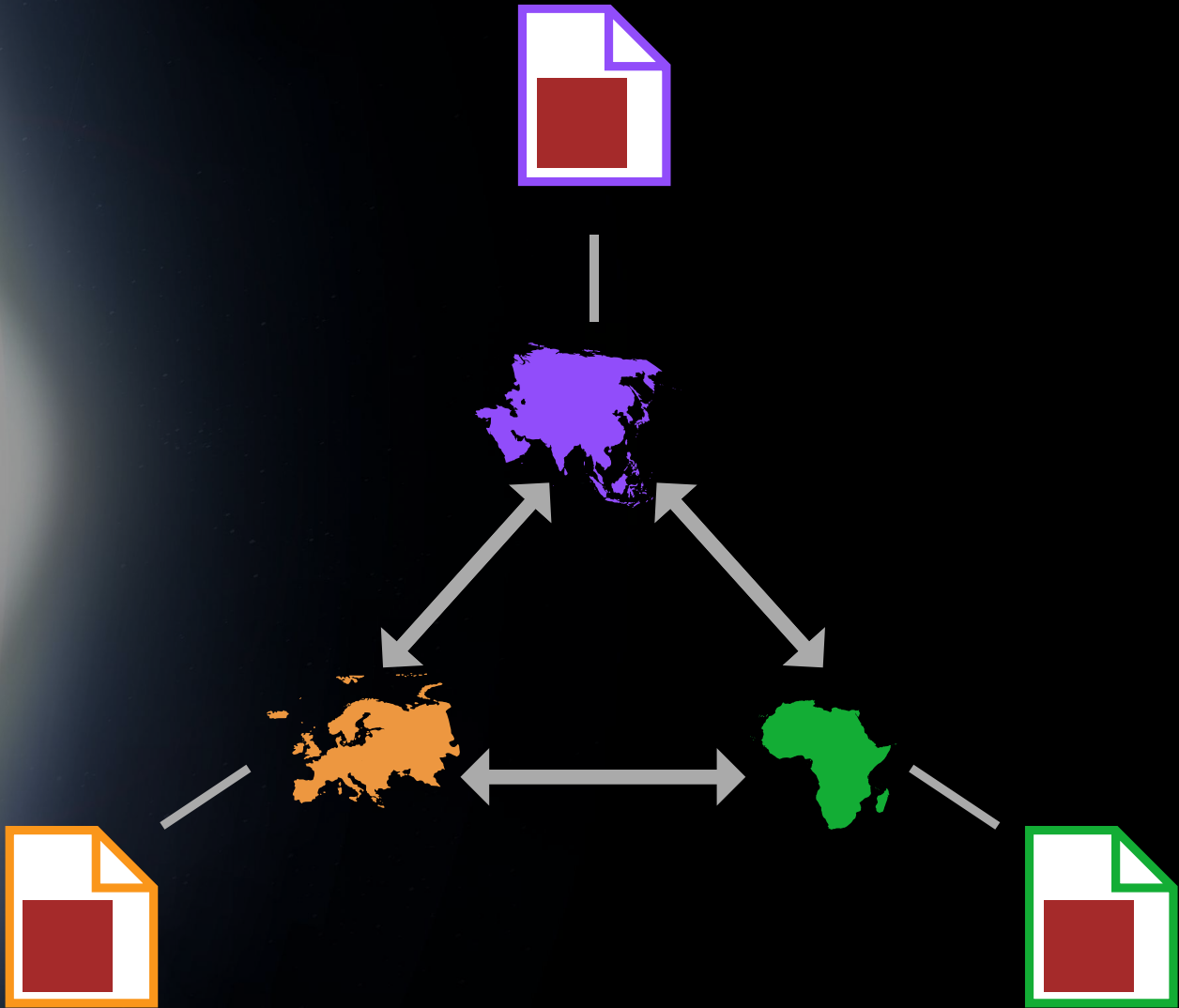
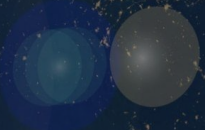
Guido Salvaneschi



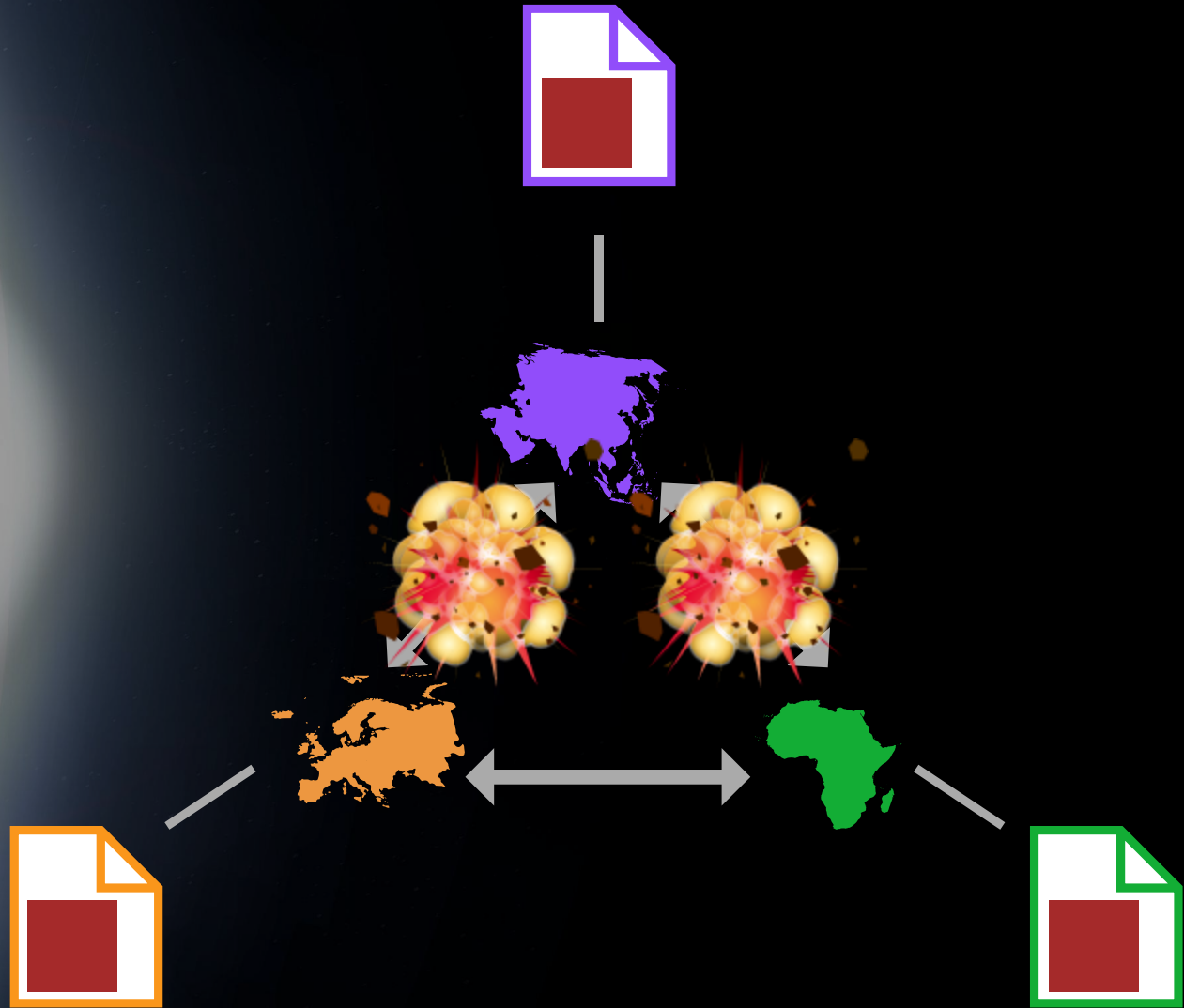
Distributed System



Distributed System: Replicated Data



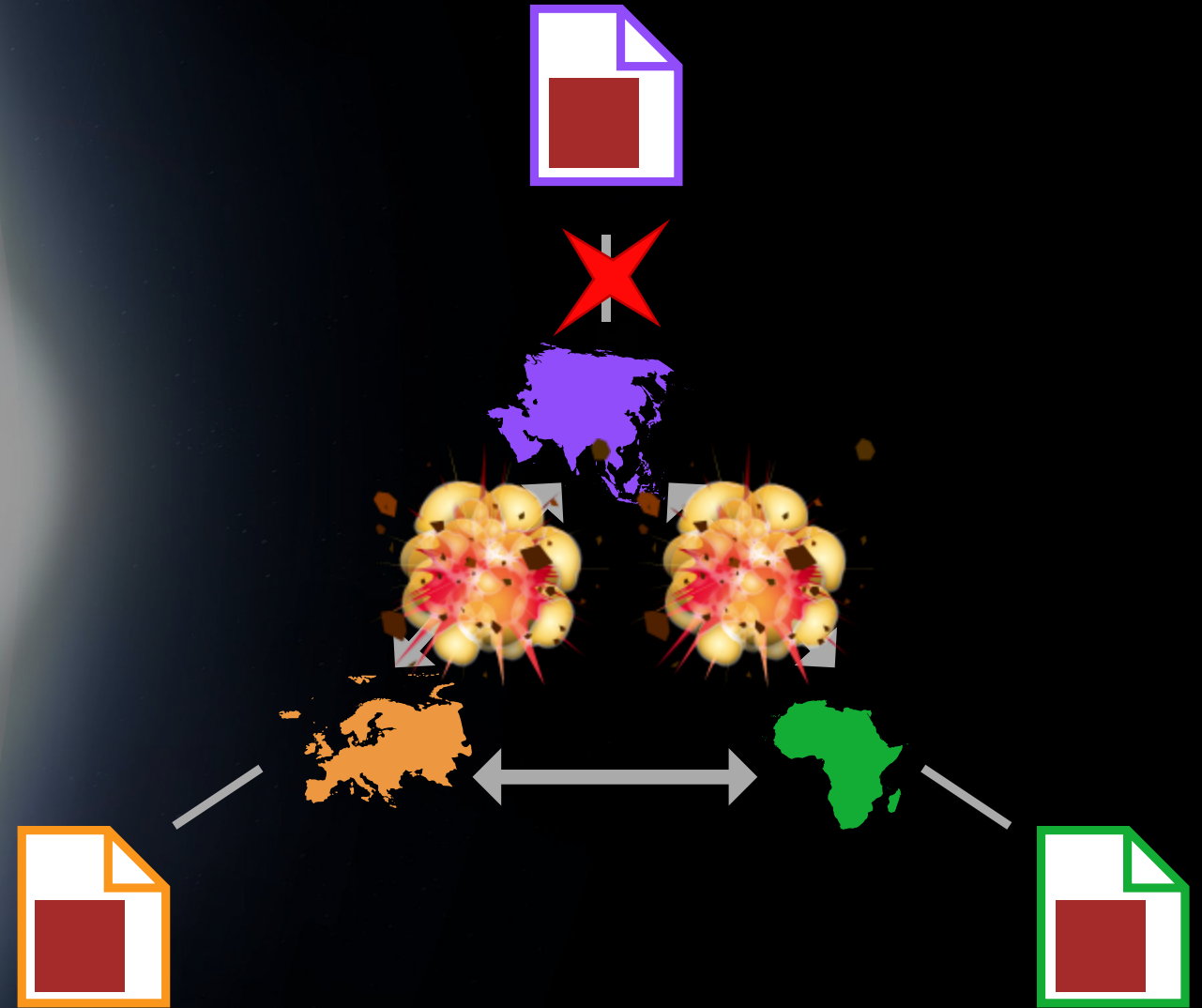
Distributed System: Replicated Data



Distributed System: Replicated Data

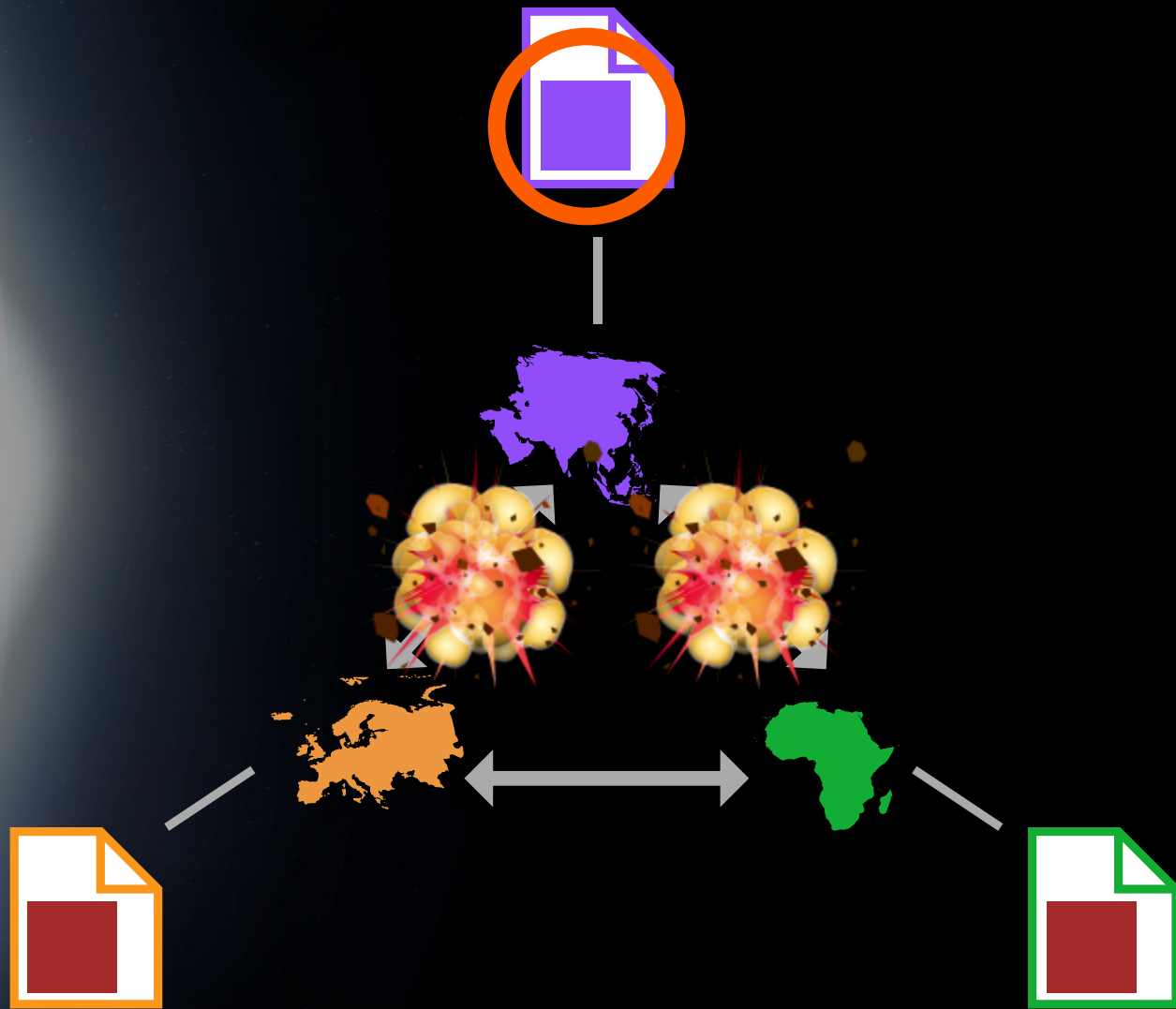
Consistency

Availability

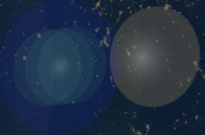


Distributed System: Replicated Data

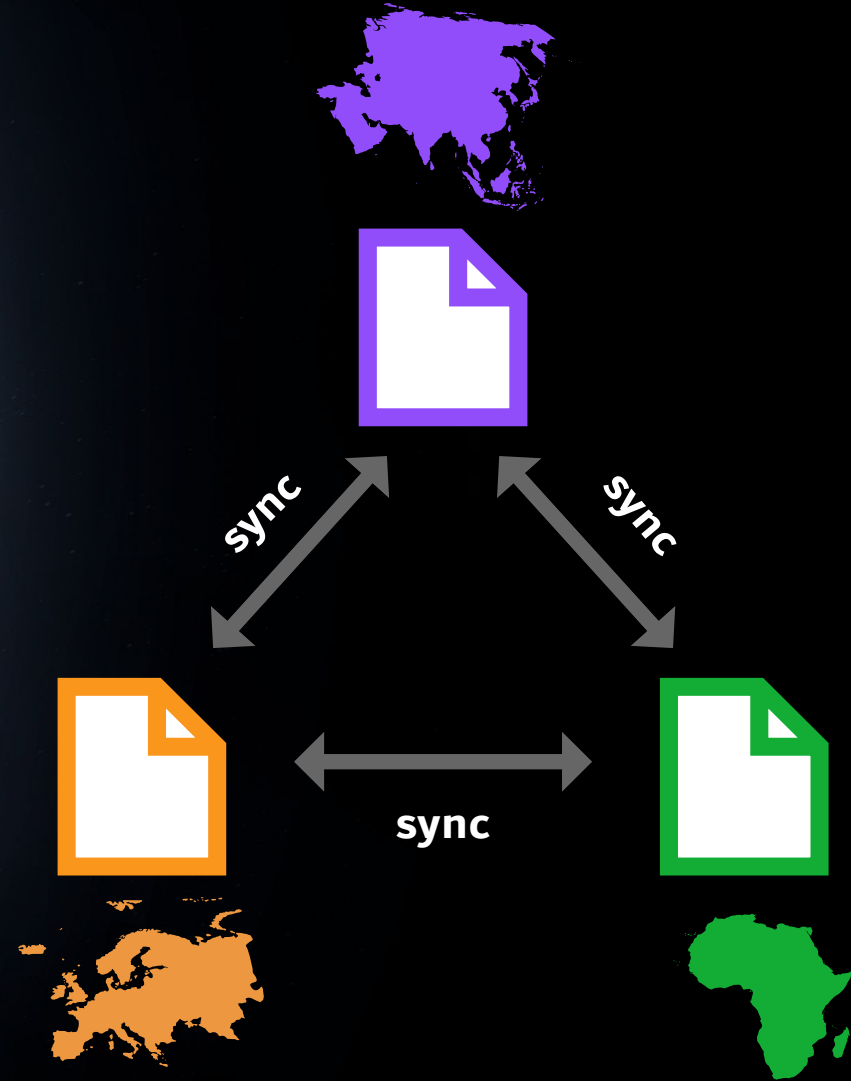
 **Consistency**
Availability



Distributed System: Replicated Data



Eventual Consistency



Conflict-Free Replicated Datatypes (CRDTs)

Conflict-Free Replicated Data Types*

Marc Shapiro^{1,5}, Nuno Preguiça^{1,2}, Carlos Baquero³, and Marek Zawirski^{1,4}

¹ INRIA, Paris, France

² CITI, Universidade Nova de Lisboa, Portugal

³ Universidade do Minho, Portugal

⁴ UPMC, Paris, France

⁵ LIP6, Paris, France

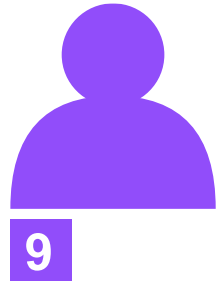
Abstract. Replicating data under Eventual Consistency (EC) allows any replica to accept updates without remote synchronisation. This ensures performance and scalability in large-scale distributed systems (e.g., clouds). However, published EC approaches are ad-hoc and error-prone. Under a formal Strong Eventual Consistency (SEC) model, we study sufficient conditions for convergence. A data type that satisfies these conditions is called a Conflict-free Replicated Data Type (CRDT). Replicas of any CRDT are guaranteed to converge in a self-stabilising manner, despite any number of failures. This paper formalises two popular approaches (state- and operation-based) and their relevant sufficient conditions. We study a number of useful CRDTs, such as sets with clean semantics, supporting both *add* and *remove* operations, and consider in depth the more complex Graph data type. CRDT types can be composed to develop large-scale distributed applications, and have interesting theoretical properties.

Keywords: Eventual Consistency, Replicated Shared Objects, Large-Scale Distributed Systems.

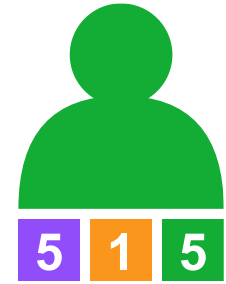
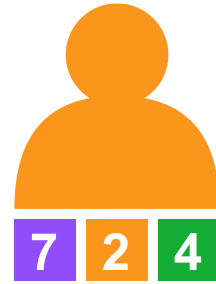
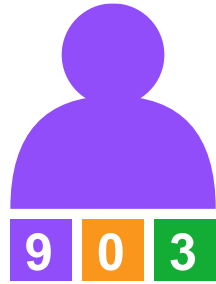
Conflict-Free Replicated Datatypes (CRDTs)



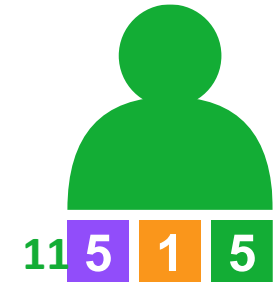
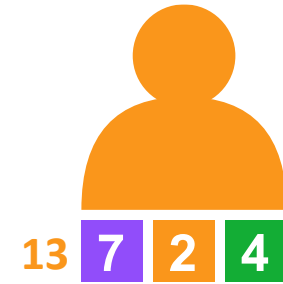
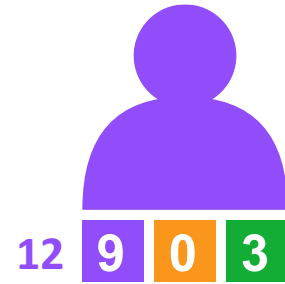
Conflict-Free Replicated Datatypes (CRDTs)



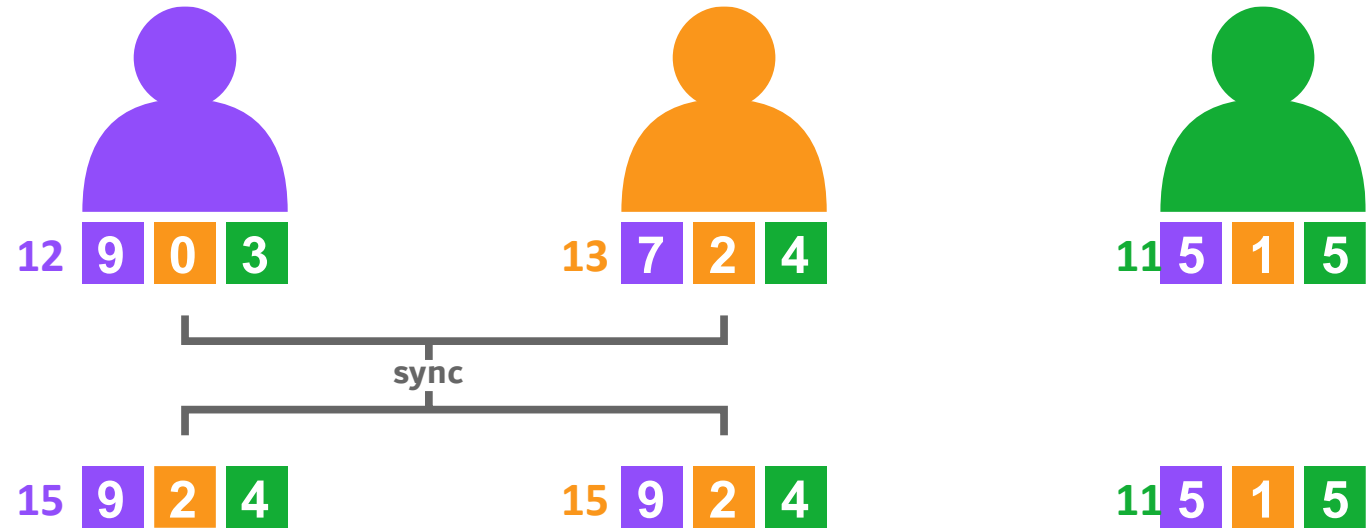
Conflict-Free Replicated Datatypes (CRDTs)



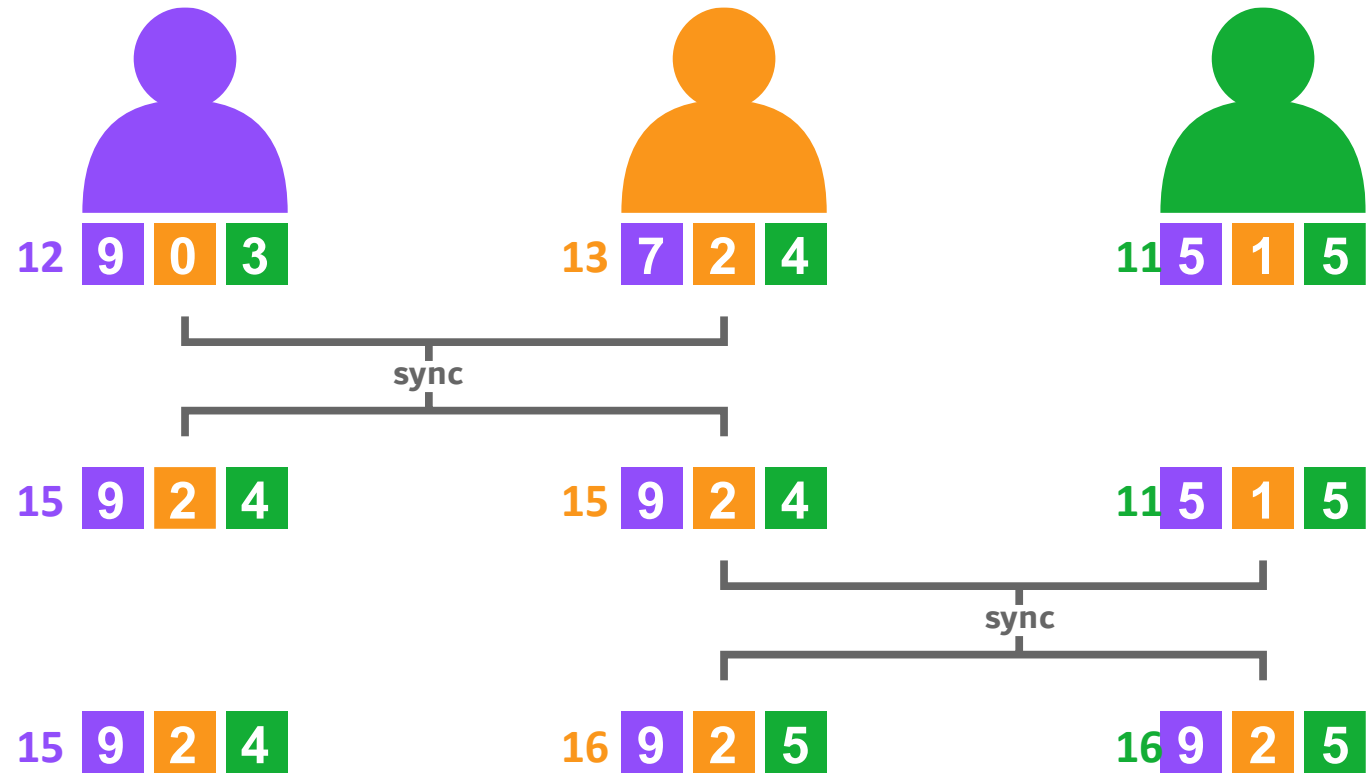
Conflict-Free Replicated Datatypes (CRDTs)



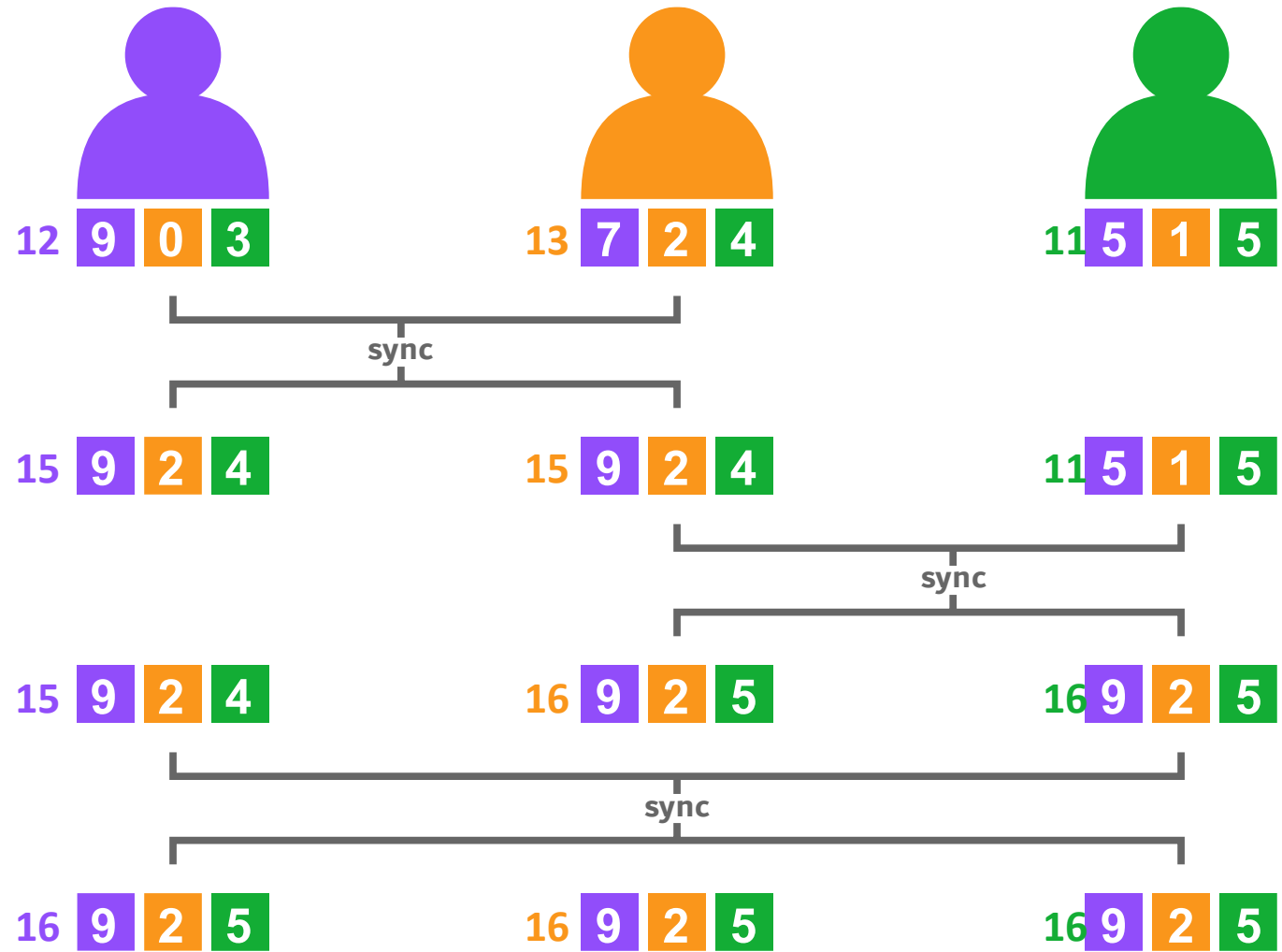
Conflict-Free Replicated Datatypes (CRDTs)



Conflict-Free Replicated Datatypes (CRDTs)



Conflict-Free Replicated Datatypes (CRDTs)



CRDTs
Guarantee
Eventual
Consistency

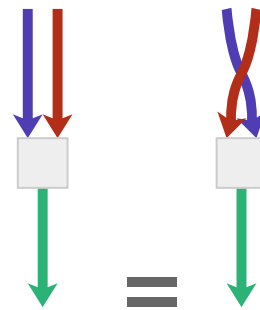
Eventual Consistency is guaranteed when syncing is:

CRDTs Guarantee Eventual Consistency

Eventual Consistency is guaranteed when syncing is:

Commutative

The order is
irrelevant



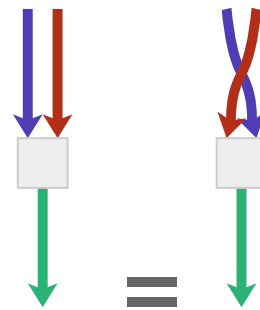
$$f(x, y) = f(y, x)$$

CRDTs Guarantee Eventual Consistency

Eventual Consistency is guaranteed when syncing is:

Commutative

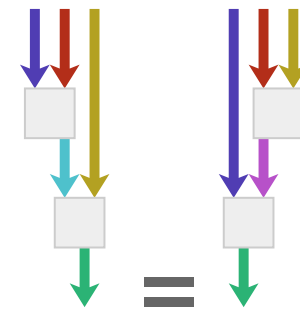
The order is irrelevant



$$f(x, y) = f(y, x)$$

Associative

The grouping is irrelevant

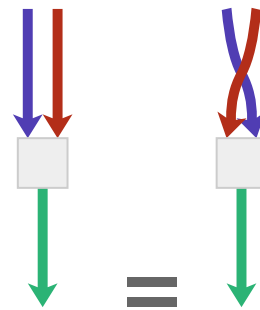


$$f(x, f(y, z)) = f(f(x, y), z)$$

CRDTs Guarantee Eventual Consistency

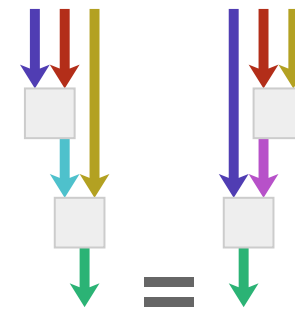
Eventual Consistency is guaranteed when syncing is:

Commutative
The order is irrelevant



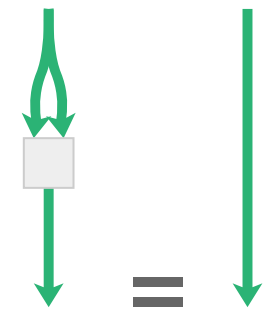
$$f(x, y) = f(y, x)$$

Associative
The grouping is irrelevant



$$f(x, f(y, z)) = f(f(x, y), z)$$

Idempotent
Syncing equals adds no new information



$$f(x, x) = x$$

CRDTs Beyond Simple Counters



Counters: Growing Counter, Positive-Negative Counter, Bounded Counters, ...

Sets: Growing Sets, Two-Phase Set, Observed-Remove Set, ...

Maps: Add-wins Observed-remove map, ...

Registers: Last-Write-Wins, ...

Building Correct CRDTs Is Hard



Martin Kleppmann @martin@nondeterministic.computer
@martinkl

Today in “distributed systems are hard”: I wrote down a simple CRDT algorithm that I thought was “obviously correct” for a course I’m teaching. Only 10 lines or so long. Found a fatal bug only after spending hours trying to prove the algorithm correct. 🤔

11:48 PM · Nov 12, 2020

38 Retweets 5 Quotes 537 Likes 65 Bookmarks



Assessing the understandability of a distributed algorithm by tweeting buggy pseudocode

Martin Kleppmann

Abstract

Designing algorithms for distributed systems has a reputation of being a difficult and error-prone task, but this difficulty is rarely measured or quantified in any way. This report tells the story of one informal experiment, in which users on Twitter were invited to identify the bug in an incorrect CRDT algorithm. Over the following 11 hours, at least 16 people (many of whom are professional software engineers) made attempts to find the bug, but most were unsuccessful. The two people who did identify the bug were both PhD students specialising in CRDTs. This result may serve as evidence of the difficulty of designing correct CRDT algorithms.



Propel

Type-Check
CRDT
Convergence

T

CRDT Type System



Propel

Type-Check
CRDT
Convergence

T

CRDT Type System

Check



Automatic Property Prover



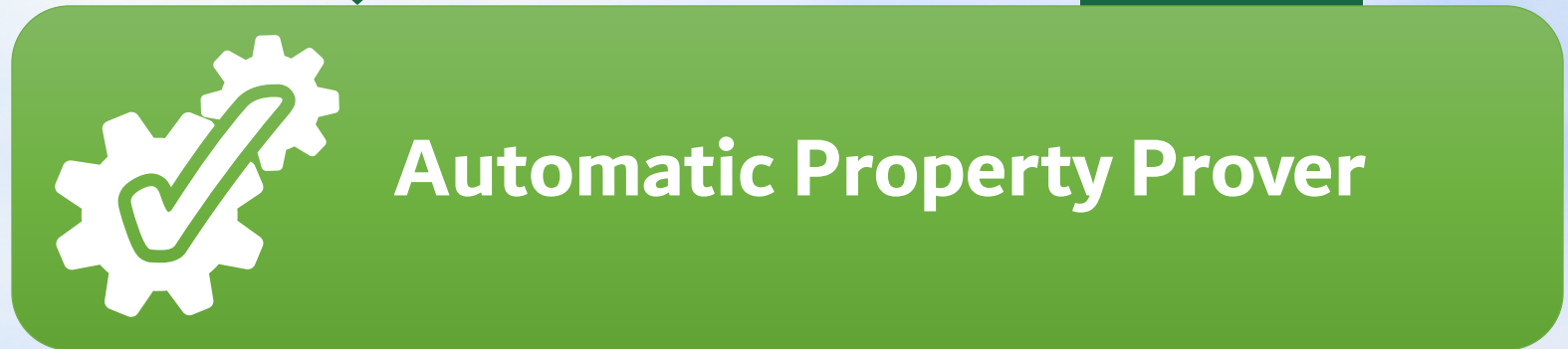
Propel

Type-Check
CRDT
Convergence



Check

Lookup





Properties in Types

```
def merge =  
  prop.rec[(Comm & Assoc & Idem) := (List[Num], List[Num]) => List[Num]]: merge =>  
    case (Nil, _) => Nil  
    case (_, Nil) => Nil  
    case (x :: xs, y :: ys) => max(y, x) :: merge(xs, ys)
```



Properties in Types

```
def merge =  
  prop.rec[(Comm & Assoc & Idem) := (List[Num], List[Num]) => List[Num]]: merge =>  
    case (Nil, _) => Nil  
    case (_, Nil) => Nil  
    case (x :: xs, y :: ys) => max(y, x) :: merge(xs, ys)
```



Properties in Types

```
def merge =  
  prop.rec[(Comm & Assoc & Idem) := (List[Num], List[Num]) => List[Num]]: merge =>  
    case (Nil, _) => Nil  
    case (_, Nil) => Nil  
    case (x :: xs, y :: ys) => max(y, x) :: merge(xs, ys)
```



Properties in Types

```
def merge =  
  prop.rec[(Comm & Assoc & Idem) := (List[Num], List[Num]) => List[Num]]: merge =>  
    case (Nil, _) => Nil  
    case (_, Nil) => Nil  
    case (x :: xs, y :: ys) => max(y, x) :: merge(xs, ys)
```

```
def max =  
  prop.rec[(Comm & Assoc & Idem) := (Num, Num) => Num]: max =>  
    ...
```




Properties in Types

```
def merge =  
  prop.rec[(Comm & Assoc & Idem) := (List[Num], List[Num]) => List[Num]]: merge =>  
    case (Nil, _) => Nil  
    case (_, Nil) => Nil  
    case (x :: xs, y :: ys) => max(y, x) :: merge(xs, ys)
```

```
def max =  
  prop.rec[(Comm & Assoc & Idem) := (Num, Num) => Num]: max =>  
    ...
```



Properties in Types Enable Composition

```
def zipWith[P >: (Comm & Assoc & Idem), T] =  
  prop.rec[(P := (T, T) => T) => (P := (List[T], List[T]) => List[T])]:  
    zipWith => f =>  
      case (Nil, y) => y  
      case (x, Nil) => x  
      case (x :: xs, y :: ys) => f(x, y) :: zipWith(f)(xs, ys)
```

Γ

Properties in Types Enable Composition

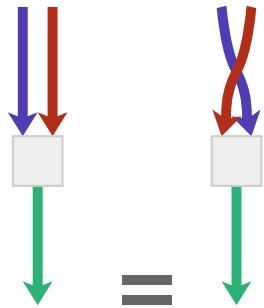
```
def zipWith[P >: (Comm & Assoc & Idem), T] =  
  prop.rec[(P := (T, T) =>: T) => (P := (List[T], List[T]) =>: List[T])]:  
    zipWith => f =>  
      case (Nil, y) => y  
      case (x, Nil) => x  
      case (x :: xs, y :: ys) => f(x, y) :: zipWith(f)(xs, ys)  
  
def merge =  
  prop[(Comm & Assoc & Idem) := (List[Num], List[Num]) =>: List[Num]]:  
    zipWith(max)
```



Properties of Binary Functions

Commutativity

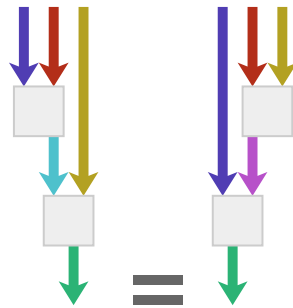
The sync order is irrelevant



$$f(x, y) = f(y, x)$$

Associativity

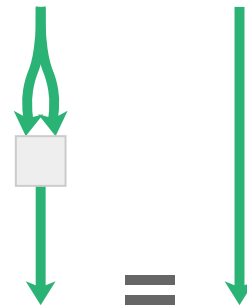
The sync order is still irrelevant



$$f(x, f(y, z)) = f(f(x, y), z)$$

Idempotency

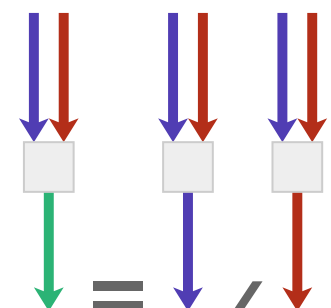
Syncing equals adds no new information



$$f(x, x) = x$$

Selectivity

A function will always return one of its argument.



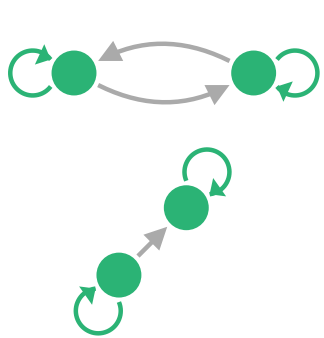
$$f(x, y) = x \vee f(x, y) = y$$

Example: projections, maximum, conditionals



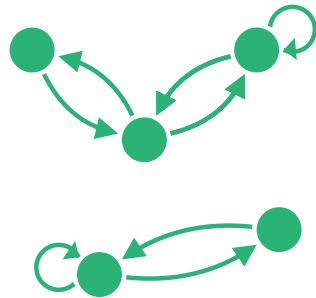
Properties of Binary Relations

Reflexivity



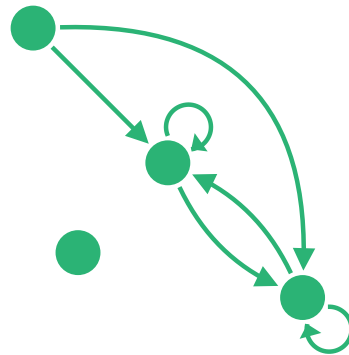
$$f(x, x) = T$$

Symmetry



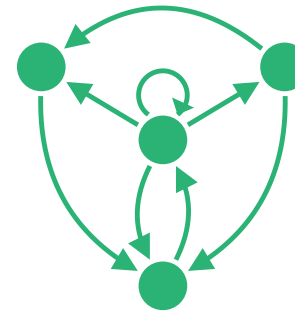
$$f(x, y) = T \\ \vdash f(y, x) = T$$

Transitivity



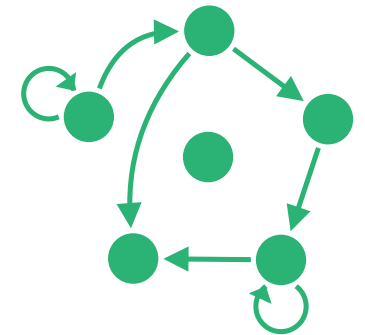
$$f(x, y) = T \wedge f(y, z) = T \\ \vdash f(x, z) = T$$

Connectivity



$$f(x, y) = T \vee f(y, x) = T$$

Antisymmetry



$$f(x, y) = T \wedge x \neq y \\ \vdash f(y, x) = \perp$$



Propel Proofs Example

```
def max = prop.rec[Comm := (BitVec, BitVec) => BitVec]: max =>
(a, b) => (a, b) match
  case (Nil, b) => b
  case (a, Nil) => a
  case (B0(x), B0(y)) => B0(max(x, y))
  case (B1(x), B1(y)) => B1(max(x, y))
  case (B0(x), B1(y)) => if equals(max(x, y), y) then B1(y) else B0(x)
  case (B1(x), B0(y)) => if equals(max(x, y), y) then B1(x) else B0(y)
```



Propel Proofs Example

$\max(a, b) = \max(b, a)$

Goal

```
def max = prop.rec[Comm := (BitVec, BitVec) => BitVec]: max =>
(a, b) => (a, b) match
  case (Nil, b) => b
  case (a, Nil) => a
  case (B0(x), B0(y)) => B0(max(x, y))
  case (B1(x), B1(y)) => B1(max(x, y))
  case (B0(x), B1(y)) => if equals(max(x, y), y) then B1(y) else B0(x)
  case (B1(x), B0(y)) => if equals(max(x, y), y) then B1(x) else B0(y)
```



Propel Proofs Example

$\max(a, b) = \max(b, a)$

Goal

$\Gamma =$

`equals: (Antisym & Sym & Trans) := (BitVec, BitVec) =>: Bool,
a: BitVec, b: BitVec`

Typing Context

```
def max = prop.rec[Comm := (BitVec, BitVec) =>: BitVec]: max =>  
  (a, b) => (a, b) match  
    case (Nil, b) => b  
    case (a, Nil) => a  
    case (B0(x), B0(y)) => B0(max(x, y))  
    case (B1(x), B1(y)) => B1(max(x, y))  
    case (B0(x), B1(y)) => if equals(max(x, y), y) then B1(y) else B0(x)  
    case (B1(x), B0(y)) => if equals(max(x, y), y) then B1(x) else B0(y)
```




Propel Proofs Example

$\max(a, b) = \max(b, a)$

Goal

$\Gamma =$

`equals: (Antisym & Sym & Trans) := (BitVec, BitVec) =>: Bool,
a: BitVec, b: BitVec`

Typing Context

Equalities

```
def max = prop.rec[Comm := (BitVec, BitVec) =>: BitVec]: max =>  
  (a, b) => (a, b) match  
    case (Nil, b) => b  
    case (a, Nil) => a  
    case (B0(x), B0(y)) => B0(max(x, y))  
    case (B1(x), B1(y)) => B1(max(x, y))  
    case (B0(x), B1(y)) => if equals(max(x, y), y) then B1(y) else B0(x)  
    case (B1(x), B0(y)) => if equals(max(x, y), y) then B1(x) else B0(y)
```



Propel Proofs Example

$\max(a, b) = \max(b, a)$

Goal

$\Gamma =$

`equals: (Antisym & Sym & Trans) := (BitVec, BitVec) =>: Bool,
a: BitVec, b: BitVec`

Typing Context

Equalities

Inequalities

```
def max = prop.rec[Comm := (BitVec, BitVec) =>: BitVec]: max =>  
  (a, b) => (a, b) match  
  case (Nil, b) => b  
  case (a, Nil) => a  
  case (B0(x), B0(y)) => B0(max(x, y))  
  case (B1(x), B1(y)) => B1(max(x, y))  
  case (B0(x), B1(y)) => if equals(max(x, y), y) then B1(y) else B0(x)  
  case (B1(x), B0(y)) => if equals(max(x, y), y) then B1(x) else B0(y)
```



Propel Proofs: Case Analysis

$\text{max}(\text{Nil}, b) = \text{max}(b, \text{Nil})$

Goal

$\Gamma = \text{max}: \text{Comm} := (\text{BitVec}, \text{BitVec}) \Rightarrow: \text{BitVec},$ Typing Context
 $\text{equals}: (\text{Antisym} \ \& \ \text{Sym} \ \& \ \text{Trans}) := (\text{BitVec}, \text{BitVec}) \Rightarrow: \text{Bool},$
 $a: \text{BitVec}, b: \text{BitVec}, a': \text{BitVec}$

$a = \text{Nil}$

Equalities

$a \neq \text{B0}(a')$
 $a \neq \text{B1}(a')$

Inequalities

```
def max = prop.rec[Comm := (BitVec, BitVec) =>: BitVec]: max =>
(a, b) => (a, b) match
  case (Nil, b) => b
  case (a, Nil) => a
  case (B0(x), B0(y)) => B0(max(x, y))
  case (B1(x), B1(y)) => B1(max(x, y))
  case (B0(x), B1(y)) => if equals(max(x, y), y) then B1(y) else B0(x)
  case (B1(x), B0(y)) => if equals(max(x, y), y) then B1(x) else B0(y)
```



Propel Proofs: Reflexivity

$b = b$

Goal

$\Gamma = \text{max}$:	Comm	$:= (\text{BitVec}, \text{BitVec}) \Rightarrow: \text{BitVec},$	Typing Context
	equals : (Antisym & Sym & Trans)	$:= (\text{BitVec}, \text{BitVec}) \Rightarrow: \text{Bool},$	
	$a: \text{BitVec}, b: \text{BitVec}, a': \text{BitVec}, b': \text{BitVec}, x: \text{BitVec}, y: \text{BitVec}$		

Equalities

Inequalities

```
def max = prop.rec[Comm := (BitVec, BitVec) =>: BitVec]: max =>
  (a, b) => (a, b) match
  case (Nil, b) => b
  case (a, Nil) => a
  case (B0(x), B0(y)) => B0(max(x, y))
  case (B1(x), B1(y)) => B1(max(x, y))
  case (B0(x), B1(y)) => if equals(max(x, y), y) then B1(y) else B0(x)
  case (B1(x), B0(y)) => if equals(max(x, y), y) then B1(x) else B0(y)
```



Propel Proofs: Next Case

$B0(\max(x, y)) = B0(\max(y, x))$

Goal

$\Gamma = \text{max}:$ **Comm** $:= (\text{BitVec}, \text{BitVec}) \Rightarrow: \text{BitVec},$ Typing Context
 equals: (**Antisym & Sym & Trans**) $:= (\text{BitVec}, \text{BitVec}) \Rightarrow: \text{Bool},$
 $a: \text{BitVec}, b: \text{BitVec}, a': \text{BitVec}, b': \text{BitVec}, x: \text{BitVec}, y: \text{BitVec}$

$a = B0(x)$
 $b = B0(y)$

Equalities

$a \neq \text{Nil}$
 $a \neq B1(a')$
 $b \neq \text{Nil}$

$b \neq B1(b')$

Inequalities

```
def max = prop.rec[Comm := (BitVec, BitVec) =>: BitVec]: max =>
(a, b) => (a, b) match
  case (Nil, b) => b
  case (a, Nil) => a
  case (B0(x), B0(y)) => B0(max(x, y))
  case (B1(x), B1(y)) => B1(max(x, y))
  case (B0(x), B1(y)) => if equals(max(x, y), y) then B1(y) else B0(x)
  case (B1(x), B0(y)) => if equals(max(x, y), y) then B1(x) else B0(y)
```



Propel Proofs: Constructor Elimination

$$B0(\max(x, y)) = B0(\max(y, x))$$

Goal

$\Gamma = \max:$ **Comm** $:= (\text{BitVec}, \text{BitVec}) \Rightarrow: \text{BitVec},$ Typing Context
 equals: (**Antisym & Sym & Trans**) $:= (\text{BitVec}, \text{BitVec}) \Rightarrow: \text{Bool},$
 $a: \text{BitVec}, b: \text{BitVec}, a': \text{BitVec}, b': \text{BitVec}, x: \text{BitVec}, y: \text{BitVec}$

$a = B0(x)$
 $b = B0(y)$

Equalities

$a \neq \text{Nil}$
 $a \neq B1(a')$
 $b \neq \text{Nil}$

$b \neq B1(b')$

Inequalities

```
def max = prop.rec[Comm := (BitVec, BitVec) =>: BitVec]: max =>
(a, b) => (a, b) match
  case (Nil, b) => b
  case (a, Nil) => a
  case (B0(x), B0(y)) => B0(max(x, y))
  case (B1(x), B1(y)) => B1(max(x, y))
  case (B0(x), B1(y)) => if equals(max(x, y), y) then B1(y) else B0(x)
  case (B1(x), B0(y)) => if equals(max(x, y), y) then B1(x) else B0(y)
```



Propel Proofs: Constructor Elimination

$\max(x, y) = \max(y, x)$

Goal

$\Gamma = \max:$ **Comm** $:= (\text{BitVec}, \text{BitVec}) \Rightarrow: \text{BitVec},$ Typing Context
 equals: (**Antisym & Sym & Trans**) $:= (\text{BitVec}, \text{BitVec}) \Rightarrow: \text{Bool},$
 $a: \text{BitVec}, b: \text{BitVec}, a': \text{BitVec}, b': \text{BitVec}, x: \text{BitVec}, y: \text{BitVec}$

$a = B0(x)$
 $b = B0(y)$

Equalities

$a \neq \text{Nil}$
 $a \neq B1(a')$
 $b \neq \text{Nil}$

$b \neq B1(b')$

Inequalities

```
def max = prop.rec[Comm := (BitVec, BitVec) =>: BitVec]: max =>
(a, b) => (a, b) match
  case (Nil, b) => b
  case (a, Nil) => a
  case (B0(x), B0(y)) => B0(max(x, y))
  case (B1(x), B1(y)) => B1(max(x, y))
  case (B0(x), B1(y)) => if equals(max(x, y), y) then B1(y) else B0(x)
  case (B1(x), B0(y)) => if equals(max(x, y), y) then B1(x) else B0(y)
```



Propel Proofs: Use Properties from Γ

$\max(x, y) = \max(y, x)$

Goal

Γ max: **Comm** := (BitVec, BitVec) =>: BitVec, Typing Context
 equals: (**Antisym & Sym & Trans**) := (BitVec, BitVec) =>: Bool,
 a: BitVec, b: BitVec, a': BitVec, b': BitVec, x: BitVec, y: BitVec

a = B0(x)
 b = B0(y)
 $\max(x, y) = \max(y, x)$

Equalities

a ≠ Nil b ≠ B1(b')
 a ≠ B1(a')
 b ≠ Nil

Inequalities

```
def max = prop.rec[Comm := (BitVec, BitVec) =>: BitVec]: max =>
(a, b) => (a, b) match
  case (Nil, b) => b
  case (a, Nil) => a
  case (B0(x), B0(y)) => B0(max(x, y))
  case (B1(x), B1(y)) => B1(max(x, y))
  case (B0(x), B1(y)) => if equals(max(x, y), y) then B1(y) else B0(x)
  case (B1(x), B0(y)) => if equals(max(x, y), y) then B1(x) else B0(y)
```




Propel Proofs: Use a Known Equality

$\max(y, x) = \max(y, x)$

Goal

```
 $\Gamma = \text{max}: \text{Comm} := (\text{BitVec}, \text{BitVec}) \Rightarrow: \text{BitVec},$  Typing Context  
 $\text{equals}: (\text{Antisym} \ \& \ \text{Sym} \ \& \ \text{Trans}) := (\text{BitVec}, \text{BitVec}) \Rightarrow: \text{Bool},$   
 $a: \text{BitVec}, b: \text{BitVec}, a': \text{BitVec}, b': \text{BitVec}, x: \text{BitVec}, y: \text{BitVec}$ 
```

$a = B0(x)$
 $b = B0(y)$
 $\max(x, y) = \max(y, x)$

Equalities

$a \neq \text{Nil}$ $b \neq B1(b')$
 $a \neq B1(a')$
 $b \neq \text{Nil}$

Inequalities

```
def max = prop.rec[Comm := (BitVec, BitVec) =>: BitVec]: max =>  
  (a, b) => (a, b) match  
    case (Nil, b) => b  
    case (a, Nil) => a  
    case (B0(x), B0(y)) => B0(max(x, y))  
    case (B1(x), B1(y)) => B1(max(x, y))  
    case (B0(x), B1(y)) => if equals(max(x, y), y) then B1(y) else B0(x)  
    case (B1(x), B0(y)) => if equals(max(x, y), y) then B1(x) else B0(y)
```



Propel Proofs: Next Case

if equals(max(x,y),y) then B1(y) else B0(x) = if equals(max(y,x),x) then B1(y) else B0(x) Goal

$\Gamma = \text{max: Comm} := (\text{BitVec}, \text{BitVec}) \Rightarrow: \text{BitVec},$ Typing Context
 $\text{equals: (Antisym \& Sym \& Trans)} := (\text{BitVec}, \text{BitVec}) \Rightarrow: \text{Bool},$
 $a: \text{BitVec}, b: \text{BitVec}, a': \text{BitVec}, b': \text{BitVec}, x: \text{BitVec}, y: \text{BitVec}$

$a = B0(x)$
 $b = B1(y)$

Equalities

$a \neq \text{Nil}$
 $a \neq B1(a')$
 $b \neq \text{Nil}$

$b \neq B0(b')$

Inequalities

```
def max = prop.rec[Comm := (BitVec, BitVec) =>: BitVec]: max =>
(a, b) => (a, b) match
  case (Nil, b) => b
  case (a, Nil) => a
  case (B0(x), B0(y)) => B0(max(x, y))
  case (B1(x), B1(y)) => B1(max(x, y))
  case (B0(x), B1(y)) => if equals(max(x, y), y) then B1(y) else B0(x)
  case (B1(x), B0(y)) => if equals(max(x, y), y) then B1(x) else B0(y)
```



Propel Proofs: “if” is a case expression

$B1(y) = \text{if equals}(\text{max}(y,x),x) \text{ then } B1(y) \text{ else } B0(x)$

Goal

$\Gamma = \text{max}:$ **Comm** $:= (\text{BitVec}, \text{BitVec}) \Rightarrow: \text{BitVec},$ **Typing Context**
 $\text{equals}:$ **(Antisym & Sym & Trans)** $:= (\text{BitVec}, \text{BitVec}) \Rightarrow: \text{Bool},$
 $a: \text{BitVec}, b: \text{BitVec}, a': \text{BitVec}, b': \text{BitVec}, x: \text{BitVec}, y: \text{BitVec}$

$a = B0(x)$
 $b = B1(y)$

Equalities

$\text{equals}(\text{max}(x,y),y) = \text{T}$

$a \neq \text{Nil}$
 $a \neq B1(a')$
 $b \neq \text{Nil}$

$b \neq B0(b')$

Inequalities

```
def max = prop.rec[Comm := (BitVec, BitVec) =>: BitVec]: max =>
(a, b) => (a, b) match
  case (Nil, b) => b
  case (a, Nil) => a
  case (B0(x), B0(y)) => B0(max(x, y))
  case (B1(x), B1(y)) => B1(max(x, y))
  case (B0(x), B1(y)) => if equals(max(x, y), y) then B1(y) else B0(x)
  case (B1(x), B0(y)) => if equals(max(x, y), y) then B1(x) else B0(y)
```



Propel Proofs: “if” is a case expression

$B1(y) = \text{if equals}(\text{max}(y,x),x) \text{ then } B1(y) \text{ else } B0(x)$

Goal

$\Gamma = \text{max: Comm} := (\text{BitVec}, \text{BitVec}) \Rightarrow: \text{BitVec},$ Typing Context
 $\text{equals: (Antisym \& Sym \& Trans)} := (\text{BitVec}, \text{BitVec}) \Rightarrow: \text{Bool},$
 $a: \text{BitVec}, b: \text{BitVec}, a': \text{BitVec}, b': \text{BitVec}, x: \text{BitVec}, y: \text{BitVec}$

Equalities

$a = B0(x)$
 $b = B1(y)$
 $\text{max}(x,y) = y \quad \text{equals}(\text{max}(x,y),y) = T$

Inequalities

$a \neq \text{Nil} \quad b \neq B0(b')$
 $a \neq B1(a')$
 $b \neq \text{Nil}$

```
def max = prop.rec[Comm := (BitVec, BitVec) =>: BitVec]: max =>
(a, b) => (a, b) match
case (Nil, b) => b
case (a, Nil) => a
case (B0(x), B0(y)) => B0(max(x, y))
case (B1(x), B1(y)) => B1(max(x, y))
case (B0(x), B1(y)) => if equals(max(x, y), y) then B1(y) else B0(x)
case (B1(x), B0(y)) => if equals(max(x, y), y) then B1(x) else B0(y)
```



Propel Proofs: Next Case

$B1(y) = B0(x)$

Goal

$\Gamma = \text{max}:$ **Comm** $:= (\text{BitVec}, \text{BitVec}) \Rightarrow: \text{BitVec},$ Typing Context
 equals: (**Antisym & Sym & Trans**) $:= (\text{BitVec}, \text{BitVec}) \Rightarrow: \text{Bool},$
 $a: \text{BitVec}, b: \text{BitVec}, a': \text{BitVec}, b': \text{BitVec}, x: \text{BitVec}, y: \text{BitVec}$

Equalities

$a = B0(x)$ $\text{max}(x, y) = x$
 $b = B1(y)$ $\text{equals}(\text{max}(x, y), x) = \perp$
 $\text{max}(x, y) = y$ $\text{equals}(\text{max}(x, y), y) = \top$

Inequalities

$a \neq \text{Nil}$ $b \neq \text{Bit0}(b')$
 $a \neq \text{Bit1}(a')$
 $b \neq \text{Nil}$

```
def max = prop.rec[Comm := (BitVec, BitVec) =>: BitVec]: max =>
(a, b) => (a, b) match
case (Nil, b) => b
case (a, Nil) => a
case (B0(x), B0(y)) => B0(max(x, y))
case (B1(x), B1(y)) => B1(max(x, y))
case (B0(x), B1(y)) => if equals(max(x, y), y) then B1(y) else B0(x)
case (B1(x), B0(y)) => if equals(max(x, y), y) then B1(x) else B0(y)
```



Propel Proofs: Stuck

$B1(y) = B0(x)$

Goal

$\Gamma = \text{max}:$ **Comm** $:= (\text{BitVec}, \text{BitVec}) \Rightarrow: \text{BitVec},$ Typing Context
 equals: (**Antisym & Sym & Trans**) $:= (\text{BitVec}, \text{BitVec}) \Rightarrow: \text{Bool},$
 $a: \text{BitVec}, b: \text{BitVec}, a': \text{BitVec}, b': \text{BitVec}, x: \text{BitVec}, y: \text{BitVec}$

Equalities

$a = B0(x)$
 $b = B1(y)$ $\text{equals}(\text{max}(x, y), x) = \perp$
 $\text{max}(x, y) = y$ $\text{equals}(\text{max}(x, y), y) = \top$

Inequalities

$a \neq \text{Nil}$ $b \neq B0(b')$
 $a \neq B1(a')$ $\text{max}(y, x) \neq x$
 $b \neq \text{Nil}$

```
def max = prop.rec[Comm := (BitVec, BitVec) =>: BitVec]: max =>
(a, b) => (a, b) match
  case (Nil, b) => b
  case (a, Nil) => a
  case (B0(x), B0(y)) => B0(max(x, y))
  case (B1(x), B1(y)) => B1(max(x, y))
  case (B0(x), B1(y)) => if equals(max(x, y), y) then B1(y) else B0(x)
  case (B1(x), B0(y)) => if equals(max(x, y), y) then B1(x) else B0(y)
```



Propel Proofs: Counterexample

The case $a=B0(x)$ and $b=B1(y)$ may not be commutative.

```
def max = prop.rec[Comm := (BitVec, BitVec) => BitVec]: max =>
(a, b) => (a, b) match
  case (Nil, b) => b
  case (a, Nil) => a
  case (B0(x), B0(y)) => B0(max(x, y))
  case (B1(x), B1(y)) => B1(max(x, y))
  case (B0(x), B1(y)) => if equals(max(x, y), y) then B1(y) else B0(x)
  case (B1(x), B0(y)) => if equals(max(x, y), y) then B1(x) else B0(y)
```



Propel Proofs: Counterexample

The case $a=B0(x)$ and $b=B1(y)$ may not be commutative.

Manual inspection: $\max(B1(B0(Z)), B0(Z)) \neq \max(B0(Z), B1(B0(Z)))$

```
def max = prop.rec[Comm := (BitVec, BitVec) => BitVec]: max =>
(a, b) => (a, b) match
  case (Nil, b) => b
  case (a, Nil) => a
  case (B0(x), B0(y)) => B0(max(x, y))
  case (B1(x), B1(y)) => B1(max(x, y))
  case (B0(x), B1(y)) => if equals(max(x, y), y) then B1(y) else B0(x)
  case (B1(x), B0(y)) => if equals(max(x, y), y) then B1(x) else B0(y)
```




Propel Proofs: Counterexample

The case $a=B0(x)$ and $b=B1(y)$ may not be commutative.

Manual inspection: $\max(B1(B0(Z)), B0(Z)) \neq \max(B0(Z), B1(B0(Z)))$

```
--- case (B1(x), B0(y)) => if equals(max(x, y), y) then B1(x) else B0(y)
+++ case (B1(x), B0(y)) => if equals(max(x, y), x) then B1(x) else B0(y)
```

```
def max = prop.rec[Comm := (BitVec, BitVec) => BitVec]: max =>
  (a, b) => (a, b) match
  case (Nil, b) => b
  case (a, Nil) => a
  case (B0(x), B0(y)) => B0(max(x, y))
  case (B1(x), B1(y)) => B1(max(x, y))
  case (B0(x), B1(y)) => if equals(max(x, y), y) then B1(y) else B0(x)
  case (B1(x), B0(y)) => if equals(max(x, y), x) then B1(x) else B0(y)
```



Propel Proofs: Contradiction

$B1(y) = B0(y)$

Goal

```

Γ = max:      Comm      := (BitVec, BitVec) =>: BitVec,      Typing Context
      equals: (Antisym & Sym & Trans) := (BitVec, BitVec) =>: Bool,
      a: BitVec, b: BitVec, a': BitVec, b': BitVec, x: BitVec, y: BitVec

```

Equalities

```

a = B0(x)
b = B1(y)   equals(max(x,y),y) = ⊥
max(x,y) = y equals(max(x,y),y) = ⊤

```

Inequalities

```

a ≠ Nil      b ≠ B0(b')
a ≠ B1(a')   max(x, y) ≠ y
b ≠ Nil

```

```

def max = prop.rec[Comm := (BitVec, BitVec) =>: BitVec]: max =>
(a, b) => (a, b) match
case (Nil, b) => b
case (a, Nil) => a
case (B0(x), B0(y)) => B0(max(x, y))
case (B1(x), B1(y)) => B1(max(x, y))
case (B0(x), B1(y)) => if equals(max(x, y), y) then B1(y) else B0(x)
case (B1(x), B0(y)) => if equals(max(x, y), x) then B1(x) else B0(y)

```

Equality Lifting: $\text{equals}(x, y) = \top \vdash x = y$

$\text{equals}(\text{max}(x, y), y) = \top$

$\text{max}(x, y) = y$



Is it okay?

Equality Lifting: $\text{equals}(x, y) = \top \vdash x = y$

$\text{equals}(\max(x, y), y) = \top$

$\max(x, y) = y$

Is it okay?



Theorem: Equality is the only antisymmetric, symmetric, and transitive relation

Equality Lifting: $\text{equals}(x, y) = \top \vdash x = y$

$\text{equals}(\text{max}(x, y), y) = \top$

$\text{max}(x, y) = y$

Is it okay?



Theorem: $\text{equals}: \text{Antisym} \ \& \ \text{Sym} \ \& \ \text{Trans} := (A, A) \Rightarrow \text{Bool}, \text{equals}(x, y) = \top$
 $\vdash x = y$

Theorem: $\text{equals}: \text{Antisym} \ \& \ \text{Sym} \ \& \ \text{Trans} := (A, A) \Rightarrow \text{Bool}, \text{equals}(x, y) = \perp$
 $\vdash x \neq y$

Equality Lifting: $\text{equals}(x, y) = \top \vdash x = y$

$\text{equals}(\text{max}(x, y), y) = \top$

$\text{max}(x, y) = y$

Is it okay?

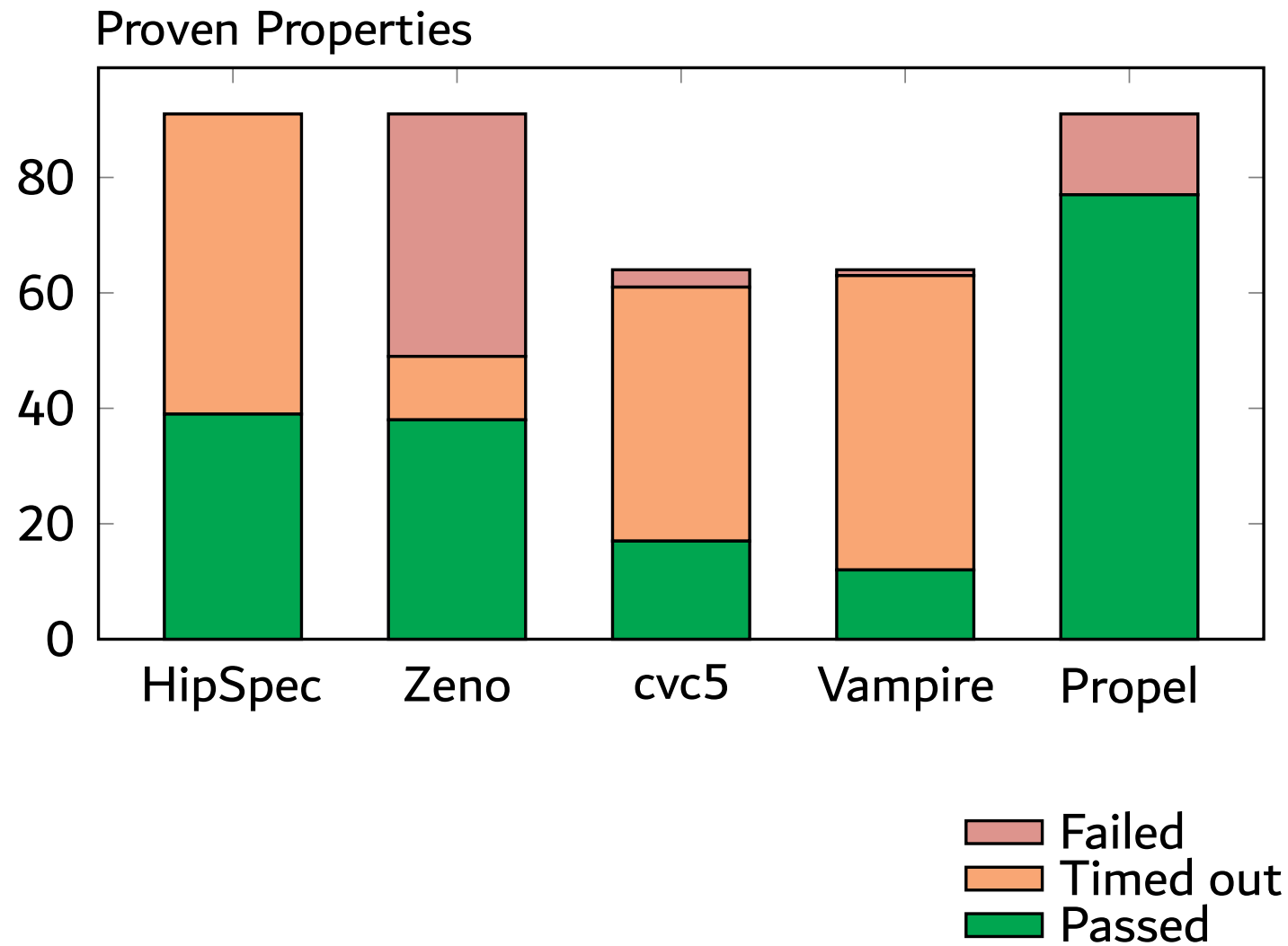


Propel discovers these theorems without any ad-hoc rules

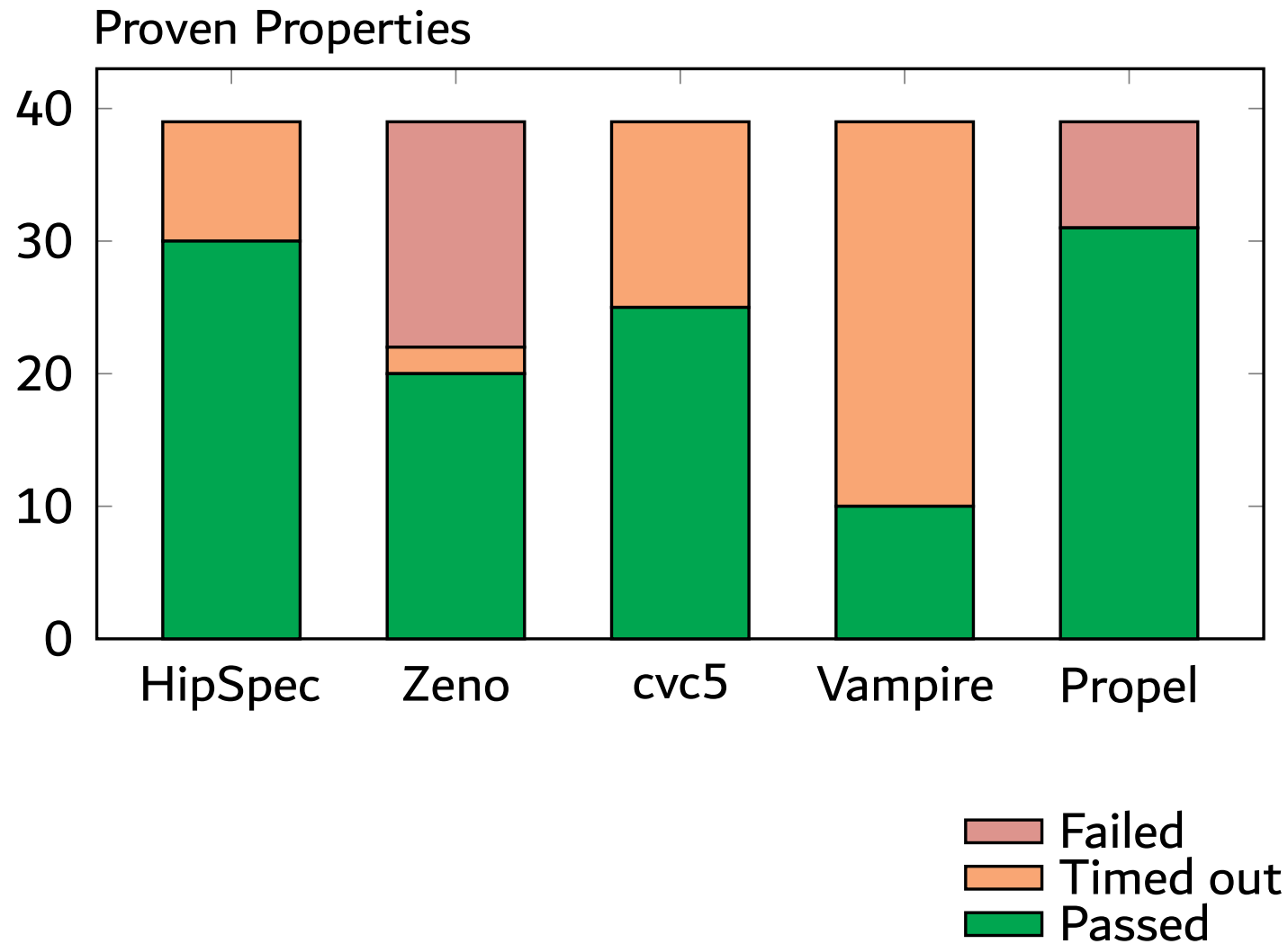
Theorem: $\text{equals}: \text{Antisym} \ \& \ \text{Sym} \ \& \ \text{Trans} := (A, A) \Rightarrow: \text{Bool}, \text{equals}(x, y) = \top$
 $\vdash x = y$

Theorem: $\text{equals}: \text{Antisym} \ \& \ \text{Sym} \ \& \ \text{Trans} := (A, A) \Rightarrow: \text{Bool}, \text{equals}(x, y) = \perp$
 $\vdash x \neq y$

Propel for CRDTs



Propel Beyond CRDTs



propel-prover.github.io



Type-check your CRDTs!

Track commutativity, associativity, idempotency,
and other algebraic properties in types

Try Propel

We compiled Propel to JavaScript using [Scala.js](#) for you to try the latest version locally in your browser.

Print reductions Print deductions Check on change

```
1 (type bool {True False})
2
3 (type Cat {Nutmeg Jake           Dovewing Socks
4 ;           |           |           |           |
5 ;           |           |           |           |
6 ; Princess Firestar Sandstorm Cloudtail
7 ;           |           |           |           |
8 ;           |           |           |           |
9 ;           |           |           |           |
10 ;           |           |           |           |
11 ;           |           |           |           |
12 ;           |           |           |           |
13 ;           |           |           |           |
14 (def cat-eq (fun Cat Cat bool)
15 (lambda (a Cat) (b Cat) (cases (Tuple a b)
16 [(Tuple Jayfeather Jayfeather) True]
17 )))
```

Checking properties for definition (rec1):

```
λ _1: (Unit). (Tuple
  (λ [trans] ancestor: Nutmeg + Jake + Dovewing + Socks + Princess
    v
    (cases (cat-parent cat) of
      Unknown → ⊥
      Tuple mother father → ((let Tuple cat-is-ancestor = rec1 U
```

Checking properties for definition (cat-is-ancestor):

```
λ [trans] ancestor: Nutmeg + Jake + Dovewing + Socks + Princess +
  v
  (cases (cat-parent cat) of
    Unknown → ⊥
    Tuple mother father → ((let Tuple cat-is-ancestor = rec1 Un
```

No decreasing recursive arguments detected for recursive definitio

Checking ... (25 s)

